



Maria Castaldo

On a Centrality Maximization Game

Joint work with:

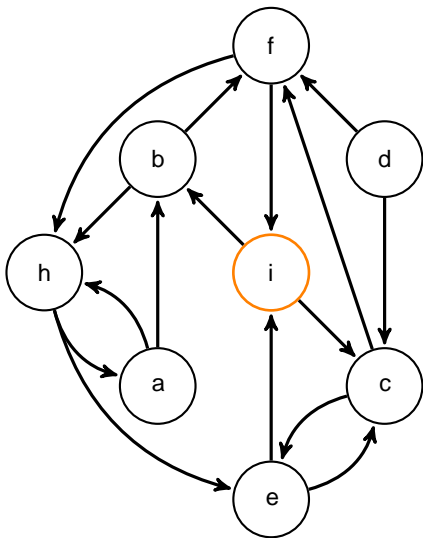
Costanza Catalano

Giacomo Como

Fabio Fagnani

Workshop on Network Dynamics in the Social, Economic, and Financial
Sciences

Bonacich Centrality



How to measure the importance i_i of a position?

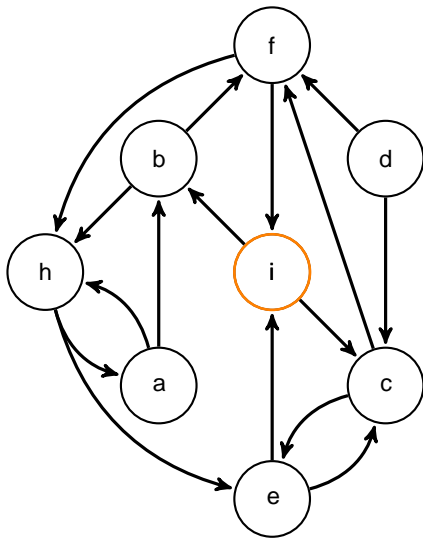
$$i_i = P^0 + (1 - \alpha) \sum_j P_{ji}^1$$

where $\sum_i P_{ii}^1 = 1$.

Bonacich centrality

Applications: visibility of a webpage on the internet.

Centrality Maximization Game

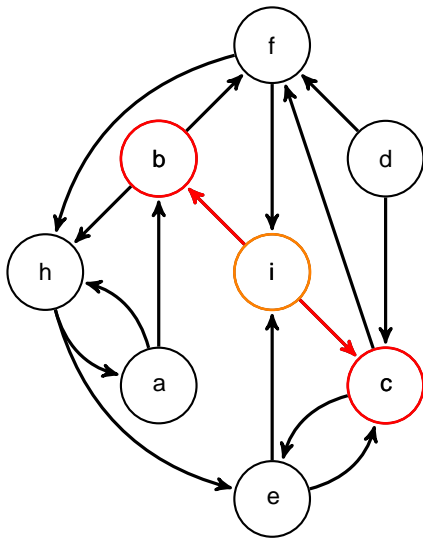


$$(V; ; ; m) = (V; fA_i g_{i2v} ; f u_i g_{i2v})$$

$V = \text{f players}$

$$x) G (x).$$

Centrality Maximization Game



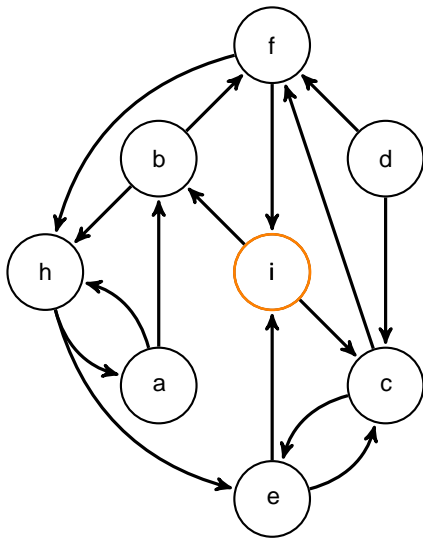
$$(V; ; ; m) = (V; fA_i; g_{i2v}; f u_i; g_{i2v})$$

$V = \{ \text{players} \}$

action $x_i \in A_i$ of node i is the set of m nodes node i points to.

$$x_i \in G(x).$$

Centrality Maximization Game



$$(V; ; ; m) = (V; fA_i g_{i2v} ; f u_i g_{i2v})$$

$V = \{ \text{players} \}$

action $x_i \in A_i$ of node i is the set of m nodes node i points to.

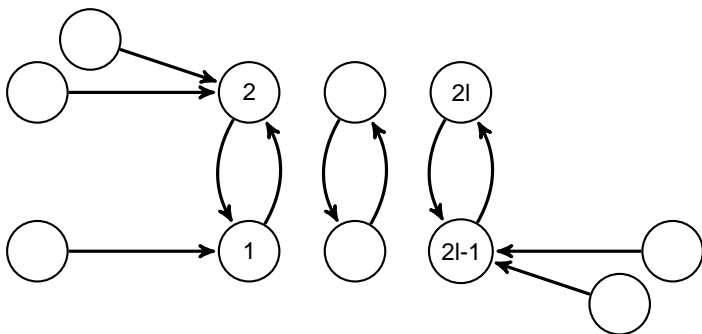
utility u_i of node i is the Bonacich centrality.

$$x \in G(x).$$

Results for $m=1$: Nash equilibria

- x is a Nash equilibrium, $G(x)$ is of type $C_2^{l;r}$, where $2 + r = |V|$:
- l number of 2-cliques;
 - r nodes with in-degree equal to zero pointing at a 2-clique.

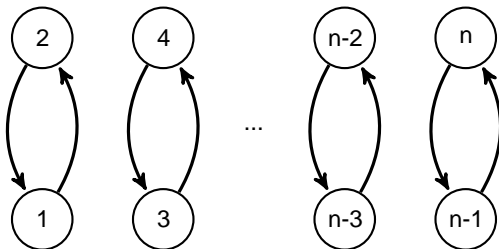
Example:



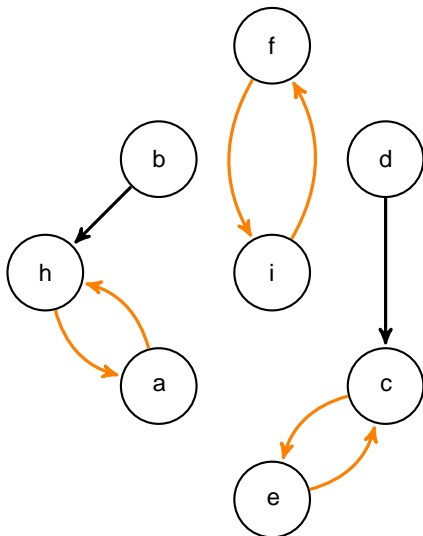
Results for $m=1$: strict Nash equilibria

If n is even, x^* is a strict Nash equilibrium, $G(x^*)$ is of type $C_2^{n-2;0}$.

If n is odd, there are no strict Nash equilibria.



Results for m=1: ordinal potential



Function : $A! N$ counting
the number of 2-cliques is an
ordinal potential ;

= 3

Results for $m=1$: Best Response Dynamics

The best response dynamics of V ; ; 1) always converges to a Nash equilibrium.

The limit set N is:

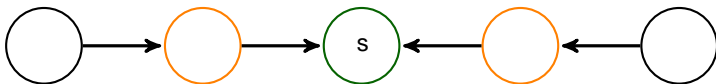
set of $C_2^{n=2;0}$,
if n is even;

set of $C_2^{(n-1)=2;1}$,
if n is odd;

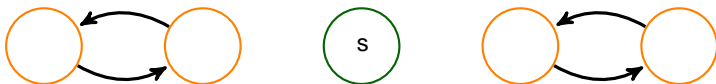
Idea of the proofs

Best Response Function

if $N_s(x) \neq \emptyset$;) $B_s(x_s) = \text{argmax}_{x \in N_s(x)} g(x)$



if $N_s(x) = \emptyset$;) $B_s(x_s) = \text{argmax}_{x \in V} g(x)$.



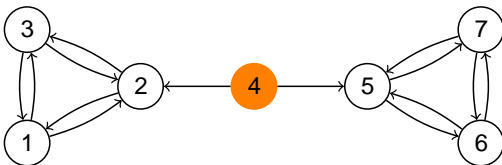
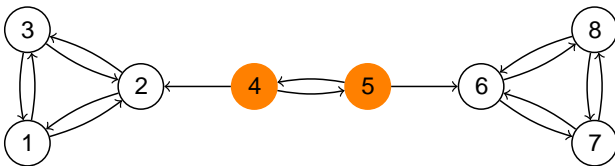
where $N_s(x)$ set of in-neighbors of s .

Results for $m=2$: Nash equilibria

If x is a Nash equilibrium:

every connected component of $G(x)$ is either a sink or a source. (if not isolated)

every source component is a **singleton** or a **2-clique**



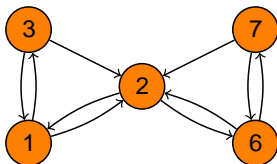
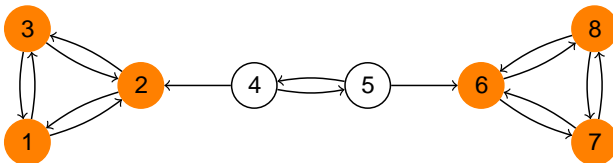
Results for $m=2$: Nash equilibria

If x is a Nash equilibrium:

every connected component of $G(x)$ is either a sink or a source. (if not isolated)

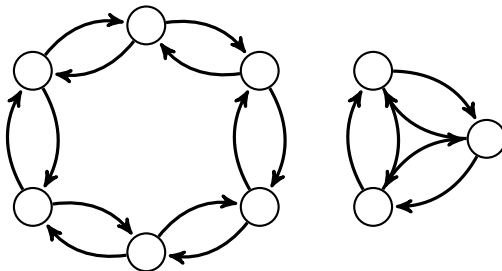
every source component is a singleton or a 2-clique

every sink component is either a cycle or the Butterfly graph



Results for $m=2$: Strict Nash equilibria

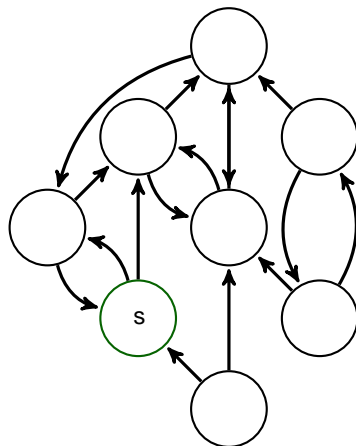
x is a strict Nash equilibrium, $G(x)$ is undirected.



Idea of the proofs

Characteristics of the best response set

if $|N_s^2(x)| \geq 2$



where $N_s^2(x) = \{j \mid \text{dist}_x(j; s) = 2\}$.

Idea of the proofs

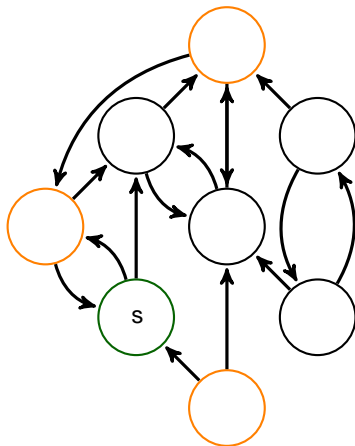
Characteristics of the best response set

if $|N_s^2(x)| \geq 2$

+

$B_s(x) \cap N_s^2(x)$

where $N_s^2(x) = \{j \mid \text{dist}_x(j; s) \leq 2\}$.



Conclusions and Next Steps

Results:

The best strategy to maximize the Bonacich centrality is to act locally.

Strict Nash equilibria are all and only undirected graphs if $m = 1$ or $m = 2$.

Existence of an ordinal potential for $m = 1$.

Next steps:

Generalization to any m .

